### **Counting the number of dominating sets of cactus chains**

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Let *G* be a simple graph of order *n*. The domination polynomial of *G* is the polynomial  $D(G, x) = \sum_{i=\gamma(G)}^{n} d(G,i)x^{i}$ , where d(G,i) is the number of dominating sets of *G* of size *i* and  $\gamma(G)$  is the domination number of *G*. The number of dominating sets of a graph *G* is D(G,1). In this paper we consider cactus chains with triangular and square blocks and study their domination polynomials.

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#### 1. Introduction

Let G = (V, E) be a simple graph. For any vertex  $v \in V(G)$ , the open neighborhood of v is the set  $N(v) = \{ u \in V(G) \mid \{u, v\} \in E(G) \}$ and the closed *neighborhood* of v is the set  $N[v] = N(v) \cup \{v\}$ . For a set  $S \subseteq V(G)$ , the open neighborhood of S is  $N(S) = \bigcup_{v \in S} N(v)$  and the closed neighborhood of S is  $N[S] = N(S) \cup S$ . A set  $S \subseteq V(G)$  is a dominating set if N[S] = V or equivalently, every vertex in  $V(G) \setminus S$  is adjacent to at least one vertex in S. The domination number  $\gamma(G)$  is the minimum cardinality of a dominating set in G. For a detailed treatment of these parameters, the reader is referred to [10]. Let D(G,i) be the family of dominating sets of a graph G with cardinality i and let d(G,i) = |D(G,i)|. The domination polynomial D(G,x)of G is defined as  $D(G, x) = \sum_{i=\gamma(G)}^{|V(G)|} d(G, i) x^i$ , where  $\gamma(G)$  is the domination number of G (see [2, 5]). Obviously, the number of dominating sets of a graph G is D(G,1) (see [4, 13]). Recently the number of the dominating sets of graph G, i.e., D(G,1) has been considered and studied in [18] with a different approach.

Domination theory have many applications in sciences and technology (see [10]). Recently the dominating set has found application in the assignment of structural domains in complex protein structures, which is an important task in bio-informatics ([8]).

We recall that the Hosoya index Z(G) of a molecule graph G, is the number of matching sets, and the Merrifield-Simmons index i(G) of graph G, is the number of independent sets. The Hosoya index of a graph has application to correlations with boiling points, entropies, calculated bond orders, as well as for coding of chemical structures. The Merrifield-Simmons index is one of the most popular topological indices in chemistry. For more information of these two indices see [1,15, 16, 19]. Note that Z(G) and i(G) can be study by the value of matching polynomial and independence polynomial at 1.

In this paper we consider a class of simple linear polymers called cactus chains. Cactus graphs were first known as Husimi trees; they appeared in the scientific literature some sixty years ago in papers by Husimi and Riddell concerned with cluster integrals in the theory of condensation in statistical mechanics [9,11,17]. We refer the reader to papers [7, 14] for some aspects of domination in cactus graphs.

A cactus graph is a connected graph in which no edge lies in more than one cycle. Consequently, each block of a cactus graph is either an edge or a cycle. If all blocks of a cactus G are cycles of the same size i, the cactus is iuniform. A triangular cactus is a graph whose blocks are triangles, i.e., a 3-uniform cactus. A vertex shared by two or more triangles is called a cut-vertex. If each triangle of a triangular cactus G has at most two cut-vertices, and each cut-vertex is shared by exactly two triangles, we say that G is a chain triangular cactus. By replacing triangles in this definitions by cycles of length 4 we obtain cacti whose every block is  $C_4$ . We call such cacti square cacti. Note that the internal squares may differ in the way they connect to their neighbors. If their cut-vertices are adjacent, we say that such a square is an ortho-square; if the cut-vertices are not adjacent, we call the square a parasquare.

In Section 2 we study the domination polynomial of the chain triangular cactus with two approach. In Section 3 we study the domination polynomials of chains of squares.

### 2. Domination polynomials of the chain triangular cactus

We call the number of triangles in G, the length of the chain. An example of a chain triangular cactus is shown in Fig. 1. Obviously, all chain triangular cacti of the same length are isomorphic. Hence, we denote the chain triangular cactus of length n by  $T_n$ . In this paper we investigate the domination polynomial of  $T_n$  by two different approach.

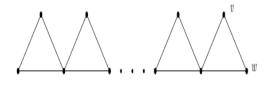


Fig. 1. The chain triangular cactus.

## **2.1** Computation of $D(T_n, x)$ using recurrence relation

In the first subsection, we use results and recurrence relations of the domination polynomial of a graph to find a recurrence relation for  $D(T_n, x)$ .

We need the following theorem:

**Theorem 1.** [5] If a graph G consists of k components  $G_1, \ldots, G_k$ , then  $D(G, x) = \prod_{i=1}^k D(G_i, x)$ .

The vertex contraction G/u of a graph G by a vertex u is the operation under which all vertices in N(u) are joined to each other and then u is deleted (see[20]).

The following theorem is useful for finding the recurrence relations for the domination polynomials of arbitrary graphs.

**Theorem 2.** [3,12] Let G be a graph. For any vertex u in G we have

$$D(G, x) = xD(G/u, x) + D(G - u, x) + xD(G - N[u], x) - (1 + x)p_u(G, x),$$

where  $p_u(G, x)$  is the polynomial counting the dominating sets of G-u which do not contain any vertex of N(u) in G.

Domination polynomial satisfies a recurrence relation for arbitrary graphs which is based on the edge and vertex elimination operations. The recurrence uses composite operations, e.g. G - e/u, which stands for (G - e)/u.

**Theorem 3.** [12] Let G be a graph. For every edge  $e = \{u, v\} \in E$ ,

$$D(G, x) = D(G - e, x) + \frac{x}{x - 1} [D(G - e/u, x) + D(G - e/v, x)]$$

$$-D(G/u, x) - D(G/v, x) - D(G - N[u], x) - D(G - N[v], x)$$

We use for graphs G = (V, E) the following vertex operation, which is commonly found in the literature. Let  $v \in V$  be a vertex of G. A vertex appending G+e (or  $G+\{v,\cdot\}$ ) denotes the graph  $(V \cup \{v'\}, E \cup \{v,v'\})$ obtained from G by adding a new vertex v' and an edge  $\{v,v'\}$  to G.

The following theorem gives recurrence relation for the domination polynomial of  $T_n$ .

**Theorem 4.** For every  $n \ge 3$ ,  $D(T_n, x) = (x^2 + 2x)D(T_{n-1}, x) + (x^2 + x)D(T_{n-2}, x)$ , with initial condition  $D(T_1, x) = x^3 + 3x^2 + 3x$  and  $D(T_2, x) = x^5 + 5x^4 + 10x^3 + 8x^2 + x$ .

**Proof.** Consider the graph  $T_n$  as shown in the following Fig. 1. Since  $T_n/u$  is isomorphic to  $T_n - u$  and  $p_u(T_n, x) = 0$ , by Theorem 2 we have:

$$\begin{split} D(T_n, x) &= x D(T_n/u, x) + D(T_n - u, x) + x D(T_n - N[u], x) \\ &- (1 + x) p_u(T_n, x) \\ &= (x + 1) D(T_{n-1} + e, x) + x D(T_{n-2} + e, x). \end{split}$$

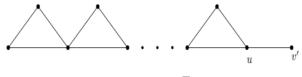


Fig. 2. The Graph  $T_{n-1} + e$ .

Note we use Theorems 1 and 2 to obtain the domination polynomial of the graph  $T_{n-1} + e$  (see Fig. 2). Suppose that v' be a vertex of degree 1 in graph  $T_{n-1} + e$  and let u be its neighbor. Note that in this case  $p_u(T_{n-1} + e, x) = 0$ . We deduce that for each  $n \in \mathbb{N}$ ,  $D(T_{n-1} + e, x) = 0$ . We deduce that for each  $n \in \mathbb{N}$ ,  $D(T_{n-1} + e, x) = x[D(T_{n-1}, x) + D(T_{n-2} + e, x) + D(T_{n-3} + e, x)]$ . Therefore  $x[D(T_n, x) = (x^2 + x)(D(T_{n-1}, x) + D(T_{n-3} + e, x)) + (x^2 + 2x)D(T_{n-2} + e, x).$  Now it's suffices to prove the following equality:  $(x^2 + x)D(T_{n-3} + e, x) + (x^2 + 2x)D(T_{n-2} + e, x) = xD(T_{n-1}, x) + (x^2 + x)D(T_{n-2}, x).$  For this purpose we use Theorem 2 for  $D(T_{n-1}, x)$ . We have

 $xD(T_{n-1}, x) = (x^2 + x)D(T_{n-2} + e, x) + x^2D(T_{n-3} + e, x).$ Now we use Theorem 2 for v' to obtain domination polynomial of  $T_{n-2} + e$ , then we have  $D(T_{n-2} + e, x) = (1 + x)D(T_{n-2}, x) + xD(T_{n-3} + e, x) - (1 + x)D(T_{n-3} + e, x).$ 

Therefore the result follows.

# **2.2** Computation of $D(T_n, x)$ by counting the number of dominating sets

In this section we shall obtain a recurrence relation for the domination polynomial of  $T_n$ . For this purpose we count the number of dominating sets of  $T_n$  with cardinality k. In other words, we first find a two variables recursive formula for  $d(T_n, k)$ .

Recently by private communication, we found that the following result also appear in [6] but were proved independently.

**Theorem 5**. The number of dominating sets of  $T_n$ with cardinality k is given by  $d(T_n, k) = 2d(T_{n-1}, k-1) + d(T_{n-1}, k-2) + d(T_{n-2}, k-1)$  $+ d(T_{n-2}, k-2).$ 

**Proof.** We shall make a dominating set of  $T_n$  with cardinality k which we denote it by  $T_n^k$ . We consider all cases:

**Case 1.** If  $T_n^k$  contains both of v and w, then we have  $T_n^k = T_{n-1}^{k-2} \cup \{v, w\}$ . In this case we have  $d(T_n, k) = d(T_{n-1}, k-2)$ .

**Case 2.** If  $\mathsf{T}_n^k$  contains only v or w (say v), then we have  $\mathsf{T}_n^k = \mathsf{T}_{n-1}^{k-1} \cup \{v\}$ . In this case we have  $d(T_n, k) = 2d(T_{n-1}, k-1)$ .

**Case 3.** If  $T_n^k$  contains none of v and w, then we can construct  $T_n^k$  by  $T_{n-2}^{k-1}$  or  $T_{n-2}^{k-2}$  as shown in Fig. 3. In this case we have  $d(T_n, k) = d(T_{n-2}, k-1) + d(T_{n-2}, k-2)$ . By adding all contributions we obtain the recurrence for  $d(T_n, k)$ .

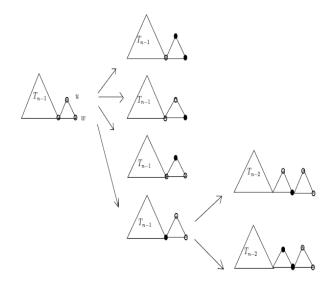


Fig. 3. Recurrence relation for  $d(T_n, k)$ .

**Corollary 1.** For every 
$$n \ge 3$$
,  
 $D(T_n, x) = (x^2 + 2x)D(T_{n-1}, x) + (x^2 + x)D(T_{n-2}, x).$ 

**Proof.** It follows from Theorem 5 and the definition of the domination polynomial.

We mention here the Hosoya index of a graph G is the total number of matchings of G and the Merrifield-Simmons index is the total number of its independent sets. Motivation by these indices, we are interested to count the total number of dominating set of a graph which is equal to D(G,1). Here we present a recurrence relation to the total number of the chain triangular cactus.

**Theorem 6**. The enumerating sequence  $\{t_n\}$  for the number of dominating sets in  $T_n$   $(n \ge 2)$  is  $t_n = 3t_{n-1} + 2t_{n-2}$ 

with initial values  $t_0 = 2$ ,  $t_1 = 7$ .

**Proof.** Since  $t_n = D(T_n, 1)$ , it follows from Corollary 1.

### Counting the number of dominating sets of chains of squares

By replacing triangles in the definitions of triangular cactus, by cycles of length 4 we obtain cacti whose every block is  $C_4$ . We call such cacti, square cacti. An example of a square cactus chain is shown in Fig. 4. We see that the internal squares may differ in the way they connect to their neighbors. If their cut-vertices are adjacent, we say that such a square is an ortho-square; if the cut-vertices are not adjacent, we call the square a para-square.

### 3.1 Domination polynomial of para-chain square cactus graphs

In this subsection we consider a para-chain of length n,  $Q_n$ , as shown in Fig. 4. We shall obtain a recurrence relation for the domination polynomial of  $Q_n$ . As usual we denote the number of dominating sets of  $Q_n$  by  $d(Q_n, k)$ . The following theorem gives a recurrence relation for  $D(Q_n, x)$ .

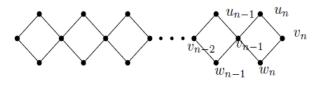


Fig. 4. Para-chain square cactus graphs.

We need the following Lemma for finding domination polynomial of the  $Q_n$ .

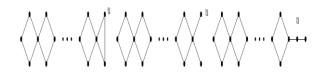


Fig. 5. Graphs  $Q_n^{\Delta}$ ,  $Q_{n'}$  and  $Q_n(2)$ , respectively.

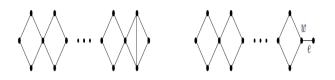


Fig. 6. Graphs  $(Q_n + e)/w$  and  $Q_n + e$ , respectively.

#### **Lemma 1.** For graphs in figures 5 and 6 have:

(*i*)  $D(Q_n^{\Delta}, x) = (1+x)D(Q_n + e, x) + xD(Q_{n-1}', x)$ , where  $D(Q_0^{\Delta}, x) = x^3 + 3x^2 + 3x$ .

(*ii*)  $D(Q_n(2), x) = x(D(Q_n + e, x) + D(Q_n, x) + D(Q_{n-1}', x))$ , where  $D(Q_0(2), x) = x^3 + 3x^2 + x$ .

(*iii*) 
$$D(Q_{n'}, x) = (1 + x)D(Q_n + e, x) - xD(Q_{n-1}', x)$$
, where  
 $D(Q_{0'}, x) = x^3 + 3x^2 + x$ .  
(*iv*)  $D(Q_n + e, x) = x(D(Q_n, x) + D(Q_{n-1}, x)) + xD(Q_{n-1}', x) + 2x^2D(Q_{n-2}', x)$   
where  $D(Q_1 + e, x) = x^5 + 5x^4 + 9x^3 + 4x^2$ .

**Proof.** The proof of parts (*i*) and (*ii*) follow from Theorems 1 and 2 for vertex *u* in graphs  $Q_n^{\Delta}$  and  $Q_n(2)$ , respectively. Note that in these cases  $p_u(G, x) = 0$ .

(*iii*) We use Theorems 1 and 2 for vertex u to obtain domination polynomial of  $Q_{n'}$ , then we have

$$D(Q_{n'}, x) = (1+x)D(Q_n + e, x) + x^2 D(Q_{n-1}', x) - (1+x)xD(Q_{n-1}', x)$$
$$= (1+x)D(Q_n + e, x) - x^2 D(Q_{n-1}', x).$$

(*iv*) We use Theorems 1 and 2 for vertex w to obtain domination polynomial of  $Q_n + e$ , as shown in Fig. 6 then  $D(Q_n + e, x) =$ 

we have  $\frac{D(Q_n + e, x) - x}{xD((Q_n + e)/w, x) + xD(Q_{n-1}', x) + xD(Q_{n-1}, x)}$ 

Now consider the graph  $(Q_n + e)/w$  as shown in Fig. 6. We use Theorems 1 and 3 for  $e = \{u, v\}$  to obtain  $D((Q_n + e)/w, x)$ , then we have

$$D((Q_n + e)/w, x) = D(Q_n, x) + \frac{x}{x-1} [D(Q_{n-1}^{\Delta}, x) + D(Q_{n-1}^{\Delta}, x) - (Q_{n-1}^{\Delta}, x) - D(Q_{n-1}^{\Delta}, x) - D(Q_{n-2}^{\Delta}, x) - D(Q_{n-2}^{\Delta}, x) + xD(Q_{n-2}^{\Delta}, x) + xD(Q_{n-2}^{\Delta}, x)] = D(Q_n, x) + 2xD(Q_{n-2}^{\Delta}, x).$$

Therefore the result follows.

**Theorem 7.** The domination polynomial of parachain  $Q_n$  is given by

$$D(Q_n, x) =$$

$$(x^3 + 2x^2 + x)D(Q_{n-1}, x) + (x^3 + 2x^2)D(Q_{n-2}, x)$$

$$+ (x^3 + 3x^2)D(Q_{n-2}', x) + (2x^4 + 4x^3)D(Q_{n-3}', x),$$

with initial conditions  $D(Q_1, x) = x^4 + 4x^3 + 6x^2$  and  $D(Q_2, x) = x^7 + 7x^6 + 21x^5 + 29x^4 + 15x^3$ .

**Proof.** Consider the labeled  $Q_n$  as shown in Figure 4. We use Theorems 1 and 2 for vertex  $u_n$  to obtain the domination polynomial of  $Q_n$ . We have

$$D(Q_n, x) = xD(Q_{n-1}^{\Delta}, x) + D(Q_{n-1}(2), x) + x^2 D(Q_{n-2}', x) - (1+x)xD(Q_{n-2}', x) = xD(Q_{n-1}^{\Delta}, x) + D(Q_{n-1}(2), x) - xD(Q_{n-2}', x).$$
(2)

Therefore by parts (i), (ii) and (iv) of Lemma 1 and equation (2) we have

$$D(Q_n, x) = x((1+x)D(Q_{n-1}+e, x) + xD(Q_{n-2}', x)) + x(D(Q_{n-1}+e, x))$$

$$\begin{aligned} &+D(Q_{n-1},x) + D(Q_{n-2}',x)) - xD(Q_{n-2}',x) \\ &= (x^2 + 2x)D(Q_{n-1} + e, x) + x^2D(Q_{n-2}', x) + xD(Q_{n-1}, x) \\ &= (x^2 + 2x)[x(D(Q_{n-1}, x) + D(Q_{n-2}, x)) + xD(Q_{n-2}', x) \\ &+ 2x^2D(Q_{n-3}', x)] + x^2D(Q_{n-2}', x) + xD(Q_{n-1}, x) \\ &= (x^3 + 2x^2 + x)D(Q_{n-1}, x) + (x^3 + 2x^2)D(Q_{n-2}, x) \\ &+ (x^3 + 3x^2)D(Q_{n-2}', x) + (2x^4 + 4x^3)D(Q_{n-3}', x). \end{aligned}$$

### **3.2 Domination polynomial of ortho-chain square cactus graphs**

In this subsection we consider a ortho-chain of length n,  $O_n$ , as shown in Fig. 7. We shall obtain a recurrence relation for the domination polynomial of  $O_n$ .

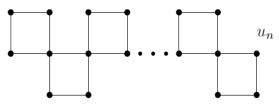


Fig. 7. Labeled ortho-chain square  $O_n$ .

We need the following Lemma for finding domination polynomial of the  $O_n$ .

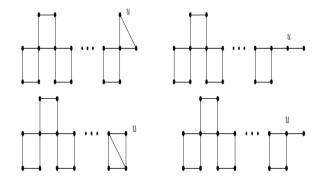


Fig. 8. Graphs  $O_n^{\Delta}$ ,  $O_n(2)$ ,  $O_{n'}$  and  $O_n + e$ , respectively.

#### **Lemma 2**. For graphs in figure 8 we have:

(i)  $D(O_n^{\Delta}, x) = (1+x)D(O_n + e, x) + xD(O_{n-1}(2), x)$ , where  $D(O_n^{\Delta}, x) = x^3 + 3x^2 + 3x$ .

(*ii*) 
$$D(O_n(2), x) = x(D(O_n + e, x) + D(O_n, x) + D(O_{n-1}(2), x))$$
  
where  $D(O_0(2), x) = x^3 + 3x^2 + x$ .

(*iii*) 
$$D(O_{n'}, x) = (1+x)D(O_n^{\Delta}, x) - xD(O_{n-1}(2), x)$$
, where

$$\begin{split} D(O_{0'}, x) &= x^4 + 4x^3 + 6x^2 + 2x \,. \\ (iv) \, D(O_n + e, x) &= x D(O_{n'}, x) + x D(O_{n-1}(2), x) + \\ x^2 D(O_{n-2}(2), x) \\ D(O_1 + e, x) &= x^5 + 5x^4 + 9x^3 + 4x^2 \,. \end{split}$$
 where

**Proof.** The proof of parts (*i*), (*ii*) and (*iv*) follow from Theorems 1 and 2 for vertex *u* in graphs  $O_n^{\Delta}$ ,  $O_n(2)$ and  $O_n + e$ , respectively. Note that in these cases  $p_u(G, x) = 0$ .

(*iii*) We use Theorems 1 and 2 for u in graphs  $O_{n'}$ . Since  $O_{n'}/u$  is isomorphic to  $O_{n'}-u$  and  $p_u(G,x) = xD(O_{n-1}(2), x)$ . So we have the result.

**Theorem 8**. The domination polynomial of parachain  $O_n$  is given by  $D(O_n, x) = xD(O_{n-1}, x) + (x^2 + 2x)D(O_{n-1} + e, x) + x^2D(O_{n-2}(2), x),$ 

with initial condition  $D(O_1, x) = x^4 + 4x^3 + 6x^2$ .

**Proof.** Consider the labeled  $O_n$  as shown in Figure 7. We use Theorems 1 and 2 for vertex  $u_n$  to obtain domination polynomial of  $O_n$ , then we have

$$\begin{split} D(O_n, x) &= x D(O_{n-1}^{\Delta}, x) + D(O_{n-1}(2), x) + \\ x^2 D(O_{n-2}(2), x) - (1+x) x D(O_{n-2}(2), x) \\ &= x D(O_{n-1}^{\Delta}, x) + D(O_{n-1}(2), x) - x D(O_{n-2}(2), x). \end{split}$$
  
Therefore by parts (i) and (ii) of Lemma 2 and this

$$\begin{split} D(O_n, x) &= x((1+x)D(O_{n-1}+e, x) + xD(O_{n-2}(2), x)) + \\ x(D(O_{n-1}+e, x) \\ &+ D(O_{n-1}, x) + D(O_{n-2}(2), x)) - xD(O_{n-2}(2), x) \\ &= (x^2 + 2x)D(O_{n-1}+e, x) + x^2D(O_{n-2}(2), x) + \\ xD(O_{n-1}, x). \end{split}$$

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equation we have

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